

Reduction-based Algorithms for Creative Telescoping

Shaoshi Chen

KLMM, Academy of Mathematics and Systems Science
Chinese Academy of Sciences

based on joint papers with A. Bostan, F. Chyzak,
H. Huang, M. Kauers, C. Koutschan,
Z. Li, M. van Hoeij, and G. Xin

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D. Zeilberger

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► Hypergeometric sums

$$\sum_{k=0}^n \frac{\binom{2k}{k}^2}{(k+1)4^{2k}} = \sum_{k=0}^n \Delta_k \left(\frac{4k \binom{2k}{k}^2}{4^{2k}} \right) = \frac{4(n+1) \binom{2n+2}{n+1}^2}{4^{2n+2}}$$

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$$L = (n+1)S_n - 4n - 2 \quad \text{and} \quad g = \frac{(2k-3n-3)k^2 \binom{n}{k}^2}{(k-n-1)^2}$$

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$$F(n) := \sum_{k=0}^n \binom{n}{k}^2 = \binom{2n}{n}$$

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- ▶ Since $f(n, k) = 0$ when $k < 0$ or $k > n$, we have

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- ▶ Taking sums on both sides of $L(f) = \Delta_k(g)$:

$$\sum_{k=-\infty}^{+\infty} L(f) = L \left(\sum_{k=-\infty}^{+\infty} f \right) = g(n, +\infty) - g(n, -\infty) = 0$$

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- ▶ Verify the initial condition:

$$F(1) = 2 = \binom{2}{1}$$

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$$\sum_{k=-a}^a (-1)^k \binom{a+b}{a+k} \binom{b+c}{b+k} \binom{c+a}{c+k} = \frac{(a+b+c)!}{a!b!c!}$$

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Hille-Hardy's identity

$$\begin{aligned} & \sum_{n=0}^{\infty} \sum_{k_1} \sum_{k_2} \frac{u^n n!}{(a+1)_n} \binom{n+a}{n-k_1} \frac{(-x)^{k_1}}{k_1!} \binom{n+a}{n-k_2} \frac{(-y)^{k_2}}{k_2!} \\ &= (1-u)^{-a-1} \exp\left\{-\frac{(x+y)^u}{1-u}\right\} \sum_n \frac{1}{n!(a+1)_n} \left(\frac{xyu}{(1-u)^2}\right)^n \end{aligned}$$

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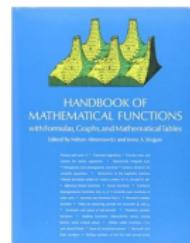
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Combinatorial
Identities

H. W. Gould



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Hypergeometric terms

A discrete function $f(n, k)$ is said to be **hypergeometric** if

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Example: $f(n, k) = \frac{(n-k)(2n+3k^2-5)}{(2k+n)(n-3k)}$

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Example: $f(n, k) = (n+k)2^n(-1)^k \frac{(n+k)!(2n-k)!(2n-2k)!}{(n+2k)!^2}$

Gosper's algorithm

In 1978, Gosper solved the telescoping problem for hypergeometric terms.

Proc. Natl. Acad. Sci. USA
Vol. 75, No. 1, pp. 40–42, January 1978
Mathematics

Decision procedure for indefinite hypergeometric summation

(algorithm/binomial coefficient identities/closed form/symbolic computation/linear recurrences)

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Example. $k! = \Delta_k(\text{No solution!})$

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Example. $\frac{(3k)!}{k!(k+1)!(k+2)!27^k} = \Delta_k(G)$

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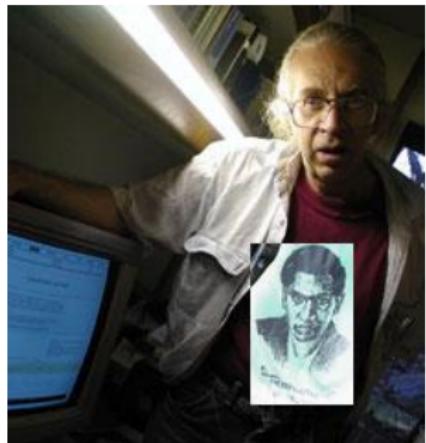
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B. Gosper

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Input: A **proper** hypergeometric term $H(n, k)$

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- ▶ Consider the hypergeometric term

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$$L_r(H) := \sum_{i=0}^r c_r H(n+i, k)$$

- ▶ Call **Gosper's algorithm** on $L_r(H)$ to check whether $\exists c_0, \dots, c_r \in K[n]$ s.t.

$$L_r(H) = \Delta_k(G_r)$$

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Input: A **proper** hypergeometric term $H(n, k)$

Output: A telescooper $L \in \mathbb{C}[n, S_n]$ s.t.

$$L(n, S_n)(H) = \Delta_k(G)$$

- ▶ Pick some $r \in \mathbb{N}$ and set $L_r = \sum_{i=0}^r c_i S_n^i$
- ▶ Consider the hypergeometric term

$$L_r(H) := \sum_{i=0}^r c_i H(n+i, k)$$

- ▶ Call **Gosper's algorithm** on $L_r(H)$ to check whether $\exists c_0, \dots, c_r \in K[n]$ s.t.

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- ▶ If all c_i 's are zero, increase r and try again

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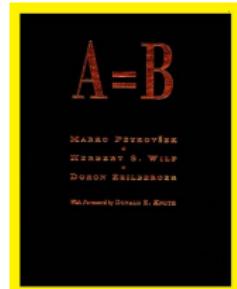
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Petkovsek, Wilf & Zeilberger

Variants of Zeilberger's algorithm

Analogous algorithms have been formulated for

- ▶ q -hypergeometric terms (Wilf-Zeilberger)
- ▶ hyperexponential terms (Almkvist-Zeilberger)
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Common feature:

Telescopers and their certificates are computed at the same time!

Telescopers

Example.

$$H = \frac{k^{10}}{n+k}$$

The telescopers of minimal order L for H is

$$L = n^{10}S_n - (n+1)^{10}$$

Telescooper

Example.

$$H = \frac{k^{10}}{n+k}$$

The telescooper of minimal order L for H is

$$L = n^{10}S_n - (n+1)^{10}$$

Guess the certificate of L ?

Certificate

$$\frac{1}{2520(n+k)} (2100k^8n^2 - 84n^3 - 68460k^6n^4 - 840n^4 - 3720n^5 + 140700k^4n^6 - 9480n^6 - 15024n^7 - 10500k^2n^8 - 14808n^8 - 8400n^9 - 79590n^2k^7 + 284235n^4k^5 - 143640n^6k^3 + 210nk^8 - 26250n^3k^6 + 133035n^5k^4 - 35700n^7k^2 + 252k^{11} + 18900k^9n - 213780k^7n^3 + 368340k^5n^5 - 110460k^3n^7 - 2100n^{10} + 1890k^9 - 1764k^7 + 1260k^5 - 378k^3 - 1260k^{10} - 294nk^2 + 700nk^4 - 588nk^6 + 63504k^{11}n^5 + 52920k^{11}n^4 + 30240k^{11}n^3 + 11340k^{11}n^2 - 2940n^2k^2 - 13080n^3k^2 - 33780n^4k^2 - 55116n^5k^2 - 57348n^6k^2 - 17360k^3n^2 - 48860k^3n^3 - 94920k^3n^4 - 135156k^3n^5 - 55440k^3n^8 - 13860k^3n^9 - 3780k^3n + 7000n^2k^4 + 31185n^3k^4 + 80850n^4k^4 + 90090n^7k^4 + 27720n^8k^4 + 57141k^5n^2 + 155610k^5n^3 + 347886k^5n^6 + 238392k^5n^7 + 110880k^5n^8 + 27720k^5n^9 + 12600k^5n - 5880n^2k^6 - 114114n^5k^6 - 123816n^6k^6 - 83160n^7k^6 - 27720n^8k^6 - 379830k^7n^4 - 469128k^7n^5 - 411840k^7n^6 - 257400k^7n^7 - 110880k^7n^8 - 27720k^7n^9 - 17640k^7n + 9405n^3k^8 + 24750n^4k^8 + 42075n^5k^8 + 47520n^6k^8 + 34650n^7k^8 + 13860n^8k^8 + 85085k^9n^2 + 398475k^9n^4 + 23100k^9n^9 + 480480k^9n^5 + 92400k^9n^8 + 235620k^9n^7 + 227150k^9n^3 + 404250k^9n^6 - 12628k^{10}n - 13860k^{10}n^9 - 152460k^{10}n^3 - 60060k^{10}n^8 - 267960k^{10}n^4 - 157080k^{10}n^7 - 271656k^{10}n^6 - 56980k^{10}n^2 - 323400k^{10}n^5 + 2520k^{11}n + 2520k^{11}n^9 + 11340k^{11}n^8 + 30240k^{11}n^7 + 52920k^{11}n^6)$$

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Very often, certificates are not needed!

Telescoping without certificates

Problem. Can we compute the **telescopers** without also computing the **certificates**?

Reduction-based Algorithms: $L(x, \partial_x)(f(x, y)) = \partial_y(g(x, y))$

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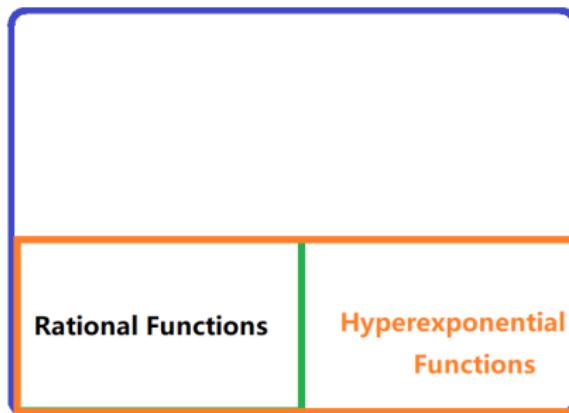


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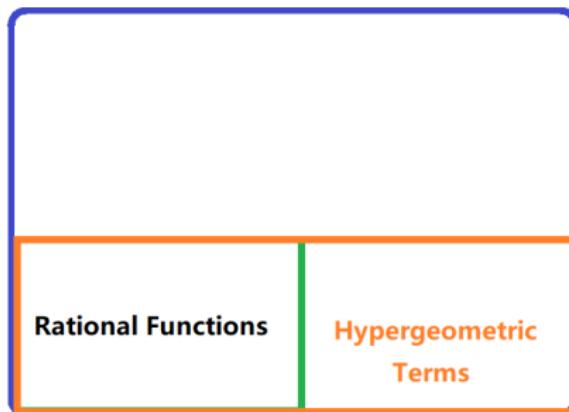


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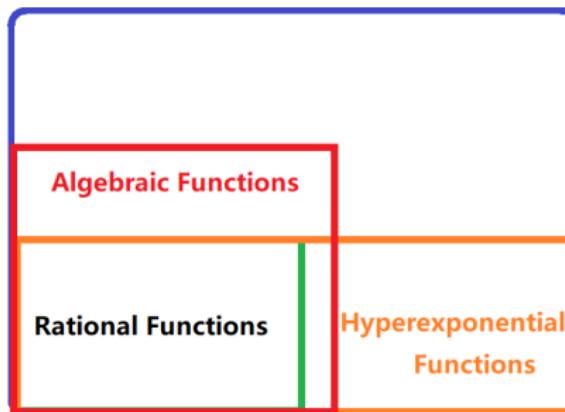


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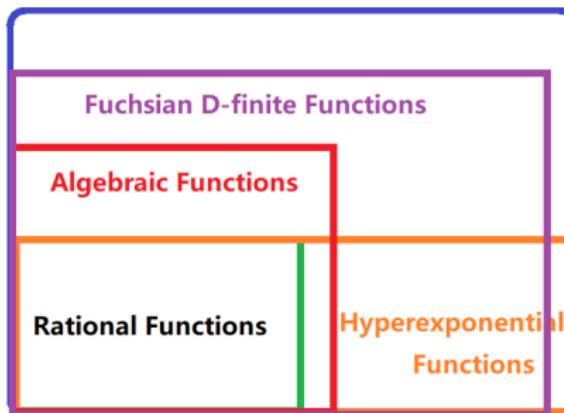


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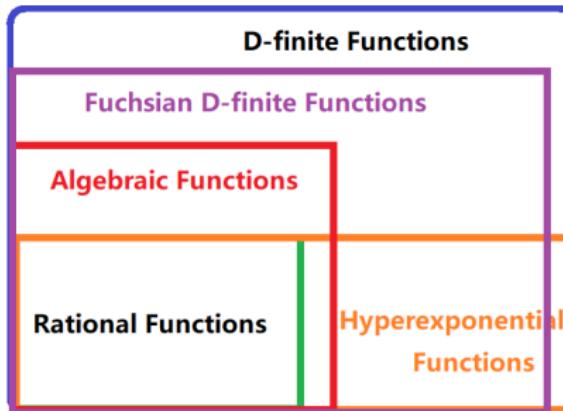


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[hal-01435877v4](https://hal.archives-ouvertes.fr/hal-01435877v4), 2017.

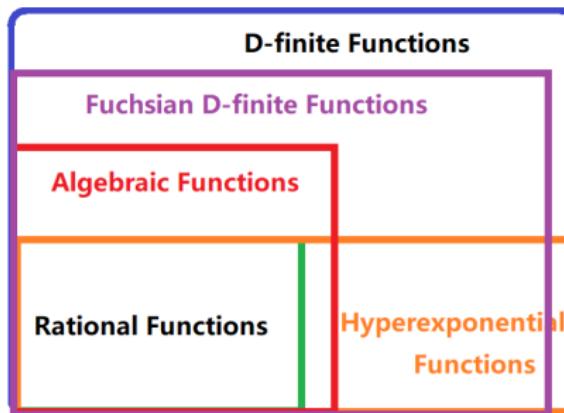


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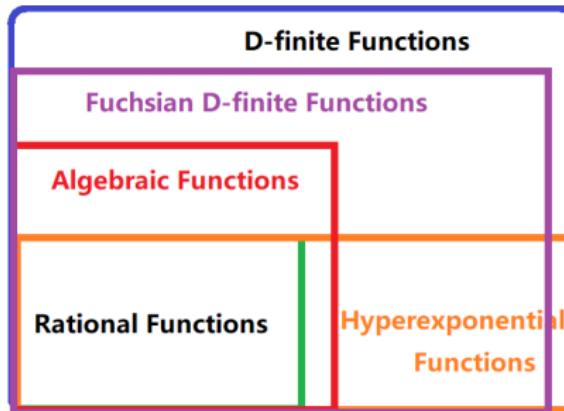


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A. Bostan, F. Chyzak, P. Lairez, B. Salvy. Generalized Hermite Reduction, Creative Telescoping and Definite Integration of D-Finite Functions. To appear in Proc. ISSAC2018.

Reduction: the rational differential case

Fact:

$$\frac{1}{(y+1)^2} + \frac{1}{y+1} = D_y \left(\frac{-1}{y+1} \right) + \frac{1}{y+1}$$

Reduction: the rational differential case

Fact:

$$\frac{1}{(y+1)^s} + \cdots + \frac{1}{y+1} = D_y \left(\frac{-(s-1)^{-1}}{(y+1)^{s-1}} + \cdots + \frac{-1}{y+1} \right) + \frac{1}{y+1}$$

Ostrogradsky-Hermite reduction: For any $f \in F(y)$,

$$f = D_y(g) + \frac{a}{b},$$

where $g \in F(y)$, $\deg_y(a) < \deg_y(b)$ and b is squarefree.

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Remark. The above reduction decomposes a rational integral as

$$\int f(y) dy = \text{rational part} + \text{transcendental part}.$$

Reduction: the rational shift case

Fact:

$$\frac{1}{k+1} = \Delta_k \left(\frac{1}{k} \right) + \frac{1}{k}$$

Reduction: the rational shift case

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Remark. The decomposition above is not unique, e.g.

$$\frac{2k+1}{k(k+1)} = \Delta_k \left(\frac{1}{k} \right) + \frac{2}{k} = \Delta_k \left(-\frac{1}{k} \right) + \frac{2}{k+1}$$

Reduction: the hyperexponential case

Let T be hyperexponential w.r.t. y with $f = D_y(T)/T \in \mathbb{F}(y)$.

$$f = \frac{D_y(r)}{r} + K \quad \rightsquigarrow \quad T = r \cdot H \quad \text{with} \quad \frac{D_y(H)}{H} = K,$$

where $K = c/d$ satisfies $\gcd(c - iD_y(d), d) = 1$ for all $i \in \mathbb{Z}$

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$$T = D_y(\dots) + \left(\frac{a}{b} + \frac{p}{d} \right) \cdot H,$$

where $a, b, p \in \mathbb{F}[y]$ with $\deg_y(a) < \deg_y(b)$, b square-free and p in a f.d. vector space V_K over \mathbb{F} .

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Reduction: the algebraic case

Let K be an algebraic extension over $\mathbb{F}(y)$ and $W := \{\omega_1, \dots, \omega_n\}$ be an integral basis of K w.r.t. y . Then any $f \in K$ is of the form

$$f = \frac{1}{D} \sum_{i=1}^n a_i \omega_i \quad \text{with } a_i, D \in \mathbb{F}[y].$$

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$$f = D_y(g) + \sum_{i=1}^n \left(\frac{p_i}{d} + \frac{q_i}{y^{\tau_i} e} \right) \omega_i, \quad \text{where}$$

- ▶ $\tau_1, \dots, \tau_n \in \mathbb{N}, d, e \in \mathbb{F}[y]$ are **squarefree**, only depend on W ;
- ▶ $p_i, q_i \in \mathbb{F}[y]$ with $\deg_y(p_i) < \deg_y(d)$ and q_i 's are in a **finitely** dimensional \mathbb{F} -vector space.

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Proposition.

$$f = D_y(g) \quad \text{for some } g \in K \quad \Leftrightarrow \quad p_i = q_i = 0$$

Telescoping via reductions

Consider

$$f = \frac{1}{xy^3 + xy + 1}$$

Find nonzero $L \in K(x)\langle D_x \rangle$ s.t. $L(f) = D_y(g)$ for $g \in K(x,y)$.

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$$D_x(f) = D_y(g_1) + \frac{-4x^2 + 6xy + 9}{x(4x^2 + 27)(xy^3 + xy + 1)}$$

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Find nonzero $L \in K(x)\langle D_x \rangle$ s.t. $L(f) = D_y(g)$ for $g \in K(x, y)$.

$$c_0(x) \cdot \frac{1}{xy^3 + xy + 1}$$

$$+ c_1(x) \cdot \frac{-4x^2 + 6xy + 9}{x(4x^2 + 27)(xy^3 + xy + 1)}$$

$$+ c_2(x) \cdot \frac{-(16x^4 + 48x^3y + 24x^2 + 81xy + 162)}{x^2(4x^2 + 27)^2(xy^3 + xy + 1)}$$

$$= 0$$

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$$\begin{aligned} G &= (4x^4 + 27x^2) \cdot g_2 + (16x^3 + 27x) \cdot g_1 + (8x^2 - 3) \cdot g_0 \\ &= \frac{-12xy^4 + 18xy^2 + 4x + 3y}{(xy^3 + xy + 1)^2} \end{aligned}$$

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Timings (in seconds)

Let

$$T = \frac{f(n, k)}{g_1(n+k)g_2(2n+k)} \frac{\Gamma(2\alpha n + k)}{\Gamma(n + \alpha k)}$$

with

- ▶ $g_i(z) = p_i(z)p_i(z+\lambda)p_i(z+\mu)$, $\alpha, \lambda, \mu \in \mathbb{N}$,
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$(m, n, \alpha, \lambda, \mu)$	Zeilberger	RCT+cert	RCT	order
(2, 0, 1, 5, 10)	354.46	58.01	4.93	4
(2, 0, 2, 5, 10)	576.31	363.25	53.15	6
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